

# Generalized weight enumerators

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November 26, 2008

# Outline

What is coding theory?

Formal definitions of codes and weights

Generalized weight enumerator

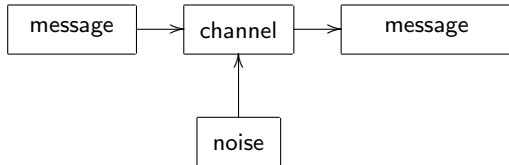
Extended weight enumerator

Matroids and the Tutte polynomial

Connections

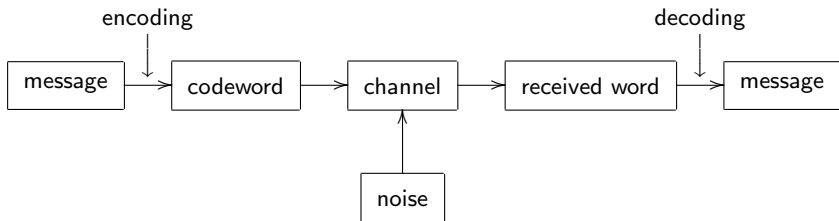
Further work

# What is coding theory?



Shannon's communication diagram

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Shannon's communication diagram

## Formal definitions of codes and weights

**Linear  $[n, k]$  code** Linear subspace  $C \subseteq \mathbb{F}_q^n$  of dimension  $k$ .  
Elements are called *(code)words*,  $n$  is called the *length*.

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### Weight enumerator

The homogeneous polynomial counting the number of words of a given weight, notation:

$$W_C(X, Y) = \sum_{w=0}^n A_w X^{n-w} Y^w.$$

## Example

The  $[7, 4]$  Hamming code over  $\mathbb{F}_2$  has generator matrix

$$G = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 \end{pmatrix}.$$

The weight enumerator is equal to

$$W_C(X, Y) = X^7 + 7X^4Y^3 + 7X^3Y^4 + Y^7.$$

## Generalized weight enumerator

For a subcode  $D \subseteq C$  we define

**Support** Union of the support of all words in  $D$ , i.e. all coordinates which are not always zero.

**Weight** Size of the support.

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## Generalized weight enumerators

The homogeneous polynomials counting for each dimension  $r = 0, \dots, k$  the number of subcodes of a given weight, notation:

$$W_C^r(X, Y) = \sum_{w=0}^n A_w^r X^{n-w} Y^w$$

## Example

The  $[7, 4]$  Hamming code has generalized weight enumerators

$$W_C^0(X, Y) = X^7$$

$$W_C^1(X, Y) = 7X^4Y^3 + 7X^3Y^4 + Y^7$$

$$W_C^2(X, Y) = 21X^2Y^5 + 7XY^6 + 7Y^7$$

$$W_C^3(X, Y) = 7XY^6 + 8Y^7$$

$$W_C^4(X, Y) = Y^7$$

## Extended weight enumerator

**Extension code** Code over some extensionfield  $\mathbb{F}_{q^m}$  having the same generator matrix as  $C$ , notation:  $C \otimes \mathbb{F}_{q^m}$ .

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### Extended weight enumerator

The polynomial “counting the number of words in an extension code”, notation:

$$W_C(X, Y, T) = \sum_{w=0}^n A_w(T) X^{n-w} Y^w.$$

Note that with  $T = q^m$  we have  $W_C(X, Y, q^m) = W_{C \otimes \mathbb{F}_{q^m}}(X, Y)$ .

## Example

The  $[7, 4]$  Hamming code has extended weight enumerator

$$\begin{aligned}W_C(X, Y, T) = & X^7 + \\ & 7(T - 1)X^4Y^3 + \\ & 7(T - 1)X^3Y^4 + \\ & 21(T - 1)(T - 2)X^2Y^5 + \\ & 7(T - 1)(T - 2)(T - 3)XY^6 + \\ & (T - 1)(T^3 - 6T^2 + 15T - 13)Y^7\end{aligned}$$



# Matroids

*Matroid theory* generalizes the notion of “linear independence”.

- Vector space: linear independent vectors, basis
- Graph: tree, minimal spanning tree
- Matroid: independent set, basis

A matroid consist of a finite set  $E$  and a set of independent sets from  $2^E$  having some defining properties.

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## Example

A code can be viewed as a matroid by considering the columns of a generator matrix and their dependance in  $\mathbb{F}_q^k$ .

## Tutte polynomial

A matroid has a *rank function*, notation  $r(A)$ , associating a non-negative integer to every subset  $A$  of  $E$ .

### Example

For matroid from a generator matrix  $G$  of a code,  $r(A)$  is the rank of the submatrix formed by the columns of  $G$  indexed by  $A$ . Furthermore,  $r(E) = k$ .

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## Tutte polynomial

The Tutte polynomial is defined by

$$t_G(X, Y) = \sum_{A \subseteq E} (X - 1)^{r(E) - r(A)} (Y - 1)^{|A| - r(A)}.$$

## Connections – some formulas

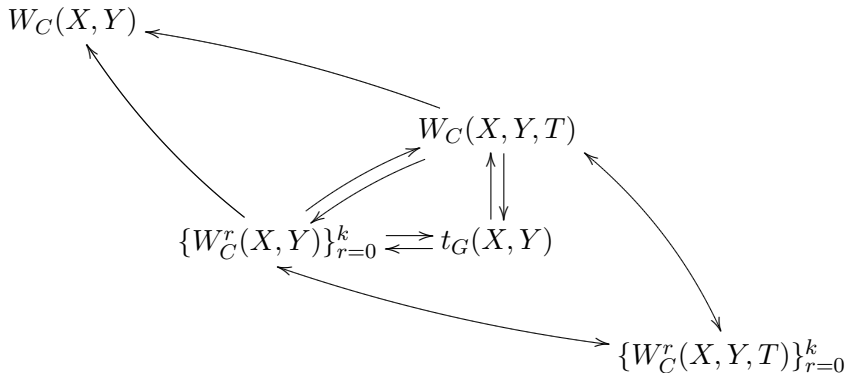
We can write the extended weight enumerator in terms of the generalized weight enumerator:

$$W_C(X, Y, T) = \sum_{r=0}^k \left( \prod_{j=0}^{r-1} (T - q^j) \right) W_C^r(X, Y);$$

and in terms of the Tutte polynomial:

$$W_C(X, Y, T) = (X - Y)^k Y^{n-k} t_G \left( \frac{X + (T - 1)Y}{X - Y}, \frac{X}{Y} \right).$$

## Connections – overview



## Further work

- Connections with other classifying polynomials:
  - Codes and zeta-function
  - Arrangement of hyperplanes and Poincaré polynomial
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  - Lattices and theta-function
- Concrete computations for special classes of codes
  - Cyclic codes
  - Algebraic geometry codes
- Characterization of extended weight enumerator

## Further work

- Extend known theory to extended weight enumerator:
  - Gleason's theory for self-dual codes
  - Codes over rings
- Probability of correct decoding
- Complexity issues



