

# Deletion and contraction in $q$ -matroids

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# $q$ -Analogues

Finite set  $\longrightarrow$  finite vectorspace over  $\mathbb{F}_q$

Example

$\binom{n}{k}$  = number of sets of size  $k$  contained in set of size  $n$

$\left[ \begin{matrix} n \\ k \end{matrix} \right]_q$  = number of  $k$ -dim subspaces of  $n$ -dim vectorspace over  $\mathbb{F}_q$

$$= \prod_{i=0}^{k-1} \frac{q^n - q^i}{q^k - q^i}$$

# $q$ -Analogues

## Example

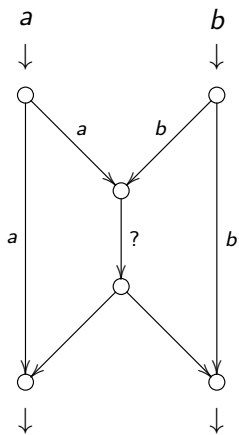
$t$ - $(v, k, \lambda)$  **design**: pair  $(X, \mathcal{B})$  with

- ▶  $X$  set with  $v$  elements (points)
- ▶  $\mathcal{B}$  family of subsets of  $X$  of size  $k$  (blocks)
- ▶ Every  $t$ -tuple of points is contained in exactly  $\lambda$  blocks

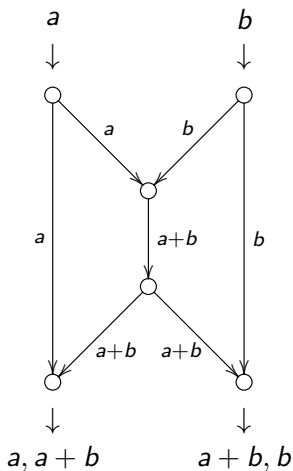
$t$ - $(v, k, \lambda; q)$   **$q$ -design**: pair  $(X, \mathcal{B})$  with

- ▶  $X$   $v$ -dim vectorspace over  $\mathbb{F}_q$
- ▶  $\mathcal{B}$  family of  $k$ -dim subspaces of  $X$  (blocks)
- ▶ Every  $t$ -dim subspace is contained in exactly  $\lambda$  blocks

# Motivation: network coding



# Motivation: network coding



Idea: send (rows of) matrices instead of vectors

Better idea: send (bases of) subspaces instead of matrices

# Motivation: network coding

Codewords are vectors:

'Ordinary' error-correcting codes

Codewords are matrices:

Rank metric codes

$q$ -analogue of 'ordinary' codes

Codewords are subspaces:

Subspace codes

Constant dimension, constant weight:  $q$ -design

# $q$ -Analogues

finite set	finite space $\mathbb{F}_q^n$
element	1-dim subspace
size	dimension
$n$	$\frac{q^n - 1}{q - 1}$
intersection	intersection
union	sum
complement	quotient space

From  $q$ -analogue to 'normal': let  $q \rightarrow 1$ .

# Matroids and $q$ -matroids

**Matroid:** a pair  $(E, \mathcal{I})$  with

- ▶  $E$  finite set;
- ▶  $\mathcal{I} \subseteq 2^E$  family of subsets of  $E$ , the *independent sets*, with:
  - (I1)  $\emptyset \in \mathcal{I}$
  - (I2) If  $A \in \mathcal{I}$  and  $B \subseteq A$  then  $B \in \mathcal{I}$ .
  - (I3) If  $A, B \in \mathcal{I}$  and  $|A| > |B|$  then there is an  $a \in A \setminus B$  such that  $B \cup \{a\} \in \mathcal{I}$ .

Examples:

- ▶ Set of vectors; independence = linear independence
- ▶ Set of edges of a graph; independence = cycle free



# Matroids and $q$ -matroids

$q$ -Matroid: a pair  $(E, \mathcal{I})$  with

- ▶  $E$  finite space;
- ▶  $\mathcal{I} \subseteq 2^E$  family of subspaces of  $E$ , the *independent spaces*, with:
  - (I1)  $\mathbf{0} \in \mathcal{I}$
  - (I2) If  $A \in \mathcal{I}$  and  $B \subseteq A$  then  $B \in \mathcal{I}$ .
  - (I3) If  $A, B \in \mathcal{I}$  and  $\dim A > \dim B$  then there is a **1-dim subspace**  $a \subseteq A$ ,  $a \not\subseteq B$  such that  $B + a \in \mathcal{I}$ .

# Why study $q$ -matroids?

Matroids generalize:

- ▶ codes
- ▶ graphs
- ▶ some designs

$q$ -Matroids generalize:

- ▶ rank metric codes
- ▶  $q$ -graphs ?
- ▶  $q$ -designs ?

## Deletion / contraction

$M = (E, \mathcal{I})$  **matroid**,  $e$  element of  $E$

Deletion  $M - e$

Ground set:  $E - \{e\}$

Independent sets: members of  $\mathcal{I}$  that do not contain  $e$

$$= \{I \in \mathcal{I} : e \notin I\}$$

Contraction  $M/e$

Ground set:  $E - \{e\}$

Independent sets: members of  $\mathcal{I}$  containing  $e$ , with  $e$  removed

$$= \{I - \{e\} : I \in \mathcal{I}, e \in I\}$$

# Deletion and contraction

$M = (E, \mathcal{I})$   $q$ -matroid,  $e$  1-dim subspace of  $E$

Deletion  $M - e$

Ground space:  $E/e$  with 'projection'  $\pi : E \rightarrow E/e$

Independent spaces: members of  $\mathcal{I}$  that do not contain  $e$ ,  
mapped to  $E/e$

$$= \{\pi(I) : I \in \mathcal{I}, e \not\subseteq I\}$$

Contraction  $M/e$

Ground space:  $E/e$  with 'projection'  $\pi : E \rightarrow E/e$

Independent spaces: members of  $\mathcal{I}$  containing  $e$ , mapped to  
 $E/e$

$$= \{\pi(I) : I \in \mathcal{I}, e \subseteq I\}$$

# Deletion and contraction

$e$  element of finite set  $E$

$\{\text{subsets not containing } e\} \cup \{\text{subsets containing } e\} = 2^E$

$$|\mathcal{I}(M - e)| + |\mathcal{I}(M/e)| = |\mathcal{I}(M)|$$

$e$  1-dimensional subspace of  $E = \mathbb{F}_q^n$

If  $A, B \in \mathcal{I}$  do not contain  $e$ , it is possible that  $\pi(A) = \pi(B)$ .

How often does that happen?

## Deletion and contraction

Total number of  $k$ -dim subspaces in  $E$ :  $\begin{bmatrix} n \\ k \end{bmatrix}_q$

Number of  $k$ -dim subspaces containing  $e$ :  $\begin{bmatrix} n-1 \\ k-1 \end{bmatrix}_q$

Number of  $k$ -dim subspaces not containing  $e$ :

$$\begin{bmatrix} n \\ k \end{bmatrix}_q - \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}_q = q^k \begin{bmatrix} n-1 \\ k \end{bmatrix}_q$$

$\implies$   $k$ -dim space in  $E/e$  has  $q^k$  pre-images in  $E$  not containing  $e$

$$\sum_{k=0}^{n-1} q^k |\{I : I \in \mathcal{I}(M - e), \dim I = k\}| + |\mathcal{I}(M/e)| = |\mathcal{I}(M)|$$

# Overview

- ▶  $q$ -Analogues attract attention nowadays because of network coding.
- ▶ We should study  $q$ -matroids for the same reasons we study matroids: they generalize several discrete structures.
- ▶ Deletion and contraction for  $q$ -matroids is defined.
- ▶ Counting on deletion and contraction in  $q$ -matroids gives extra factors  $q^k$ . This needs better notation.
- ▶ To do: deletion and contraction for the  $q$ -rank generating function and  $q$ -Tutte polynomial.
- ▶ To do: link with puncturing/shortening in rank metric codes, derived/residual  $q$ -design.

Thank you for your attention.