The *q*-analogue of matroids and their Tutte polynomial

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q-Analogues

Finite set \longrightarrow finite dimensional vectorspace over \mathbb{F}_q

Example

$$\binom{n}{k}$$
 = number of sets of size k contained in set of size n

$$\begin{bmatrix} n \\ k \end{bmatrix}_q$$
 = number of k-dim subspaces of n-dim vectorspace over \mathbb{F}_q

$$= \prod_{i=0}^{k-1} \frac{q^n - q^i}{q^k - q^i}$$

q-Analogues

$$\begin{array}{cccc} \text{finite set} & \text{finite space } \mathbb{F}_q^n \\ \text{element} & 1\text{-dim subspace} \\ \text{size} & \text{dimension} \\ n & \frac{q^n-1}{q-1} \\ \text{intersection} & \text{intersection} \\ \text{union} & \text{sum} \\ \text{complement} & \text{it depends} \end{array}$$

From *q*-analogue to 'normal': let $q \rightarrow 1$.

Candidates for complement A^c of $A \subseteq \mathbb{F}_q^n$:

- ► All vectors outside *A*But: not a space
- ► Orthogonal complement

 But: $A \cap A^{\perp}$ can be nontrivial
- Quotient space \mathbb{F}_q^n/A
- But: changes ambient space
- ▶ Subspace such that $A \oplus A^c = \mathbb{F}_q^n$ But: not unique

Matroid: a pair (E, \mathcal{I}) with

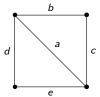
- ► *E* finite set;
- ▶ $\mathcal{I} \subseteq 2^E$ family of subsets of *E*, the *independent sets*, with:
 - (I1) $\emptyset \in \mathcal{I}$
 - (I2) If $A \in \mathcal{I}$ and $B \subseteq A$ then $B \in \mathcal{I}$.
 - (13) If $A, B \in \mathcal{I}$ and |A| > |B| then there is an $a \in A \setminus B$ such that $B \cup \{a\} \in \mathcal{I}$.

Examples:

- ► Set of vectors; independence = linear independence
- ► Set of edges of a graph; independence = cycle free

$$\left(\begin{array}{cccccc}
0 & 0 & 0 & 1 & 1 \\
0 & 1 & 1 & 0 & 0 \\
1 & 1 & 0 & 1 & 0
\end{array}\right)$$

Example



But: most matroids don't come from a matrix or graph.

A matroid is also a pair (E, r) with

(r2) If $A \subseteq B$ then r(A) < r(B).

►
$$r: 2^E \to \mathbb{N}_0$$
 a function, the *rank function*, with for all

$$A, B \in E$$
:
$$(r1) \ 0 \le r(A) \le |A|$$

(r3) $r(A \cup B) + r(A \cap B) < r(A) + r(B)$ (semimodular)

r(A) = size of largest independent set contained in A \mathcal{I} = {subsets whose size is equal to their rank}

$$\left(\begin{array}{cccccc}
0 & 0 & 0 & 1 & 1 \\
0 & 1 & 1 & 0 & 0 \\
1 & 1 & 0 & 1 & 0
\end{array}\right)$$

Example



Basis: maximal independent set wrt inclusion

Loop: element that is in no independent set (i.e., r(x) = 0)

Rank of M: rank of the ground set E

All bases have the same size.

Representable: matroid that comes form a matrix

q-Matroid: a pair (E, r) with

- ► E finite dimensional vector space;

▶
$$r: \{\text{subspaces of } E\} \rightarrow \mathbb{N}_0 \text{ a function, the } rank function, with }$$

(r2) If $A \subseteq B$ then $r(A) \le r(B)$.

(r3) $r(A+B) + r(A \cap B) < r(A) + r(B)$ (semimodular)

- for all $A, B \subseteq E$:
- $(r1) 0 \le r(A) \le \dim A$

Theorem (J. & Pellikaan, 2016)

Every \mathbb{F}_{q^m} -linear rank metric code gives a q-matroid.

Proof.

Let $E = \mathbb{F}_q^n$ and G be a generator matrix of the code. Let $A \subseteq E$ and Y a matrix whose columns span A.

Then r(A) = rk(GY) satisfies the axioms (r1),(r2),(r3).

Lemma

Matrix representation is equivalent under

- row operations over \mathbb{F}_{q^m} ;
- ightharpoonup column operations over \mathbb{F}_q .

We call a q-matroid that comes from a code representable.

Conjecture (J. & Torielli, 2017)

All *q*-matroids are representable.

That means: a q-matroid over $E = \mathbb{F}_q^n$ of rank k can be represented by a $k \times n$ matrix over a suitably large extension field \mathbb{F}_{q^m} .

Motivating evidence:

- uniform matroids are representable;
- ► the matrix has entries in an extension field.

A q-matroid could also be a pair (E, \mathcal{I}) with

- ► E finite dimensional vector space;
- \triangleright I family of subspaces of E, the *independent spaces*, with:
- (I1) $\mathbf{0} \in \mathcal{I}$. (I2) If $J \in \mathcal{I}$ and $I \subseteq J$, then $I \in \mathcal{I}$.
 - (13) If $I, J \in \mathcal{I}$ with dim $I < \dim J$, then there is some
 - 1-dimensional subspace $x \subseteq J$, $x \not\subseteq I$ with $I + x \in \mathcal{I}$.

r(A) = dimension of largest independent space contained in A $\mathcal{I} = \{$ subspaces whose dimension is equal to their rank $\}$

Let
$$E = \mathbb{F}_2^4$$
 and $\mathcal{I} = \left\{ \left\langle \begin{array}{ccc} 1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right\rangle$ and all its subspaces $\right\}$.

Let
$$E = \mathbb{F}_2^4$$
 and $\mathcal{I} = \left\{ \left\langle \begin{array}{ccc} 1 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right\rangle$ and all its subspaces $\right\}$.

$$\mathcal{I}$$
 satisfies (I1),(I2),(I3), and r satisfies (r1),(r2). But:

 $A = \left\langle \begin{array}{cccc} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{array} \right\rangle \quad B = \left\langle \begin{array}{cccc} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right\rangle$

Then $r(A+B) + r(A \cap B) = 2+1 > 1+1 = r(A) + r(B)$!

$$\mathcal{I}$$
 satisfies (I1),(I2),(I3), and r satisfies (r1),(r2). But:

$$\mathcal{I}$$
 satisfies (I1),(I2),(I3), and r satisfies (r1),(r2). But:

Problem: $(r1),(r2),(r3) \Rightarrow (l1),(l2),(l3)$; but not \Leftarrow .

Solution: find an extra axiom (I4) for $\ensuremath{\mathcal{I}}$

Lemma
Loops come in subspaces.

Corollary

If an axiom set is invariant under embedding E in a bigger space, it can not be a full axiom set for \mathcal{I} .

Theorem

A q-matroid is a pair (E, \mathcal{I}) with

- ► E finite dimensional vector space;
- $ightharpoonup \mathcal{I}$ family of subspaces of E, the independent spaces, with:
- (I1) $\mathcal{I} \neq \emptyset$. (I2) If $J \in \mathcal{I}$ and $I \subseteq J$, then $I \in \mathcal{I}$.
 - (13) If $I, J \in \mathcal{I}$ with dim $I < \dim J$, then there is some
 - 1-dimensional subspace $x \subseteq J$, $x \not\subseteq I$ with $I + x \in \mathcal{I}$.
 - (I4) Let $A, B \subseteq E$ and let I, J be maximal independent subspaces of A and B, respectively. Then there is a maximal independent subspace of A+B that is contained in I+J.

Theorem

A g-matroid is a pair (E, B) with

- ► E finite dimensional vector space;
- ▶ B family of subspaces of E, the bases, with:
 - (B1) $\mathcal{B} \neq \emptyset$ (B2) If $\mathcal{B}_1 \in \mathcal{B}_2$ and $\mathcal{B}_2 \subseteq \mathcal{B}_2$ then $\mathcal{B}_2 = \mathcal{B}_3$
 - (B2) If $B_1, B_2 \in \mathcal{B}$ and $B_1 \subseteq B_2$, then $B_1 = B_2$. (B3) If $B_1, B_2 \in \mathcal{B}$, then for every codimension 1 subspace A of B_1
 - with $B_1 \cap B_2 \subseteq B$, then for every codimension 1 subspace A of B with $A + y \in B$.
 - (B4) Let $A, B \subseteq E$ and let I, J be maximal intersections of some bases with A and B, respectively. Then there is a maximal intersection of a basis and A+B that is contained in I+J.

Duality

Let
$$r^*(A) = \dim A - r(M) + r(A^{\perp}).$$

Theorem

$$M^* = (E, r^*)$$
 is a q-matroid, i.e., r^* satisfies (r1),(r2),(r3).

Lemma

 $\mathcal{B}(M^*)$ are the orthogonal complements of bases of M.

Restriction

Let H a hyperplane in E and let $r_{M|_H}(A) = r_M(A)$.

Theorem $M|_{H} = (H, r_{M|_{H}})$ is a q-matroid, i.e., $r_{M|_{H}}$ satisfies (r1), (r2), (r3).

Lemma $\mathcal{I}(M|_H)$ are the independent spaces of M that are contained in H.

Contraction

Let e 1-dim subspace of E, $\pi: E \to E/e$ projection, A in E/e and B in E such that $e \subseteq B$ and $\pi(B) = A$. Let $r_{M/e}(A) = r_M(B) - 1$.

Theorem

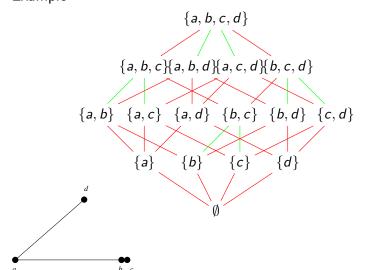
$$M/e = (E/e, r_{M/e})$$
 is a q-matroid, i.e., $r_{M/e}$ satisfies (r1),(r2),(r3).

Lemma

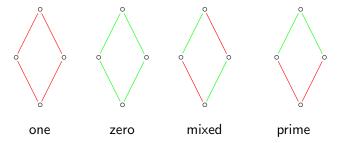
 $\mathcal{I}(M/e)$ are the independent spaces of M that contain e, projected to E/e.

Restriction, contraction, and duality

Theorem Restriction and contraction are dual operations: $(M/e)^* = M^*|_{e^{\perp}}$ and $(M|_{e^{\perp}})^* = M^*/e$.

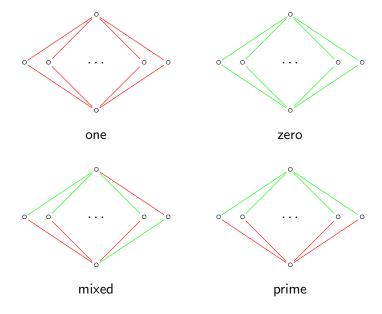


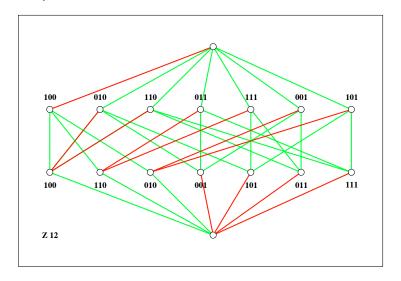
Matriod \iff only the following diamonds:

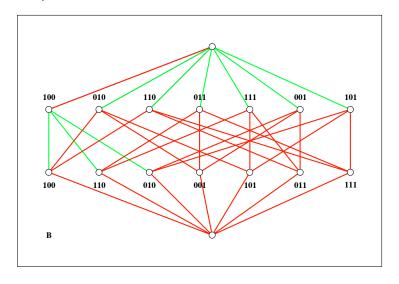


q-analogue: change Boolean lattice to subspace lattice (or another complemented modular lattice)

q-Matriod \iff only the following "diamonds":







Rank generating polynomial:

$$R(x,y) = \sum_{A \subseteq E} x^{r(M)-r(A)} y^{\dim(A)-r(A)}$$

Tutte polynomial:

classical:
$$x \to x - 1$$
, $y \to y - 1$

q: something similar but with powers of q??

Original Tutte polynomial:

$$T(x,y) = \sum_{B \in \mathcal{B}} x^{i(B)} y^{e(B)}$$

Internal/external activity uses ordering on elements of the matroid.

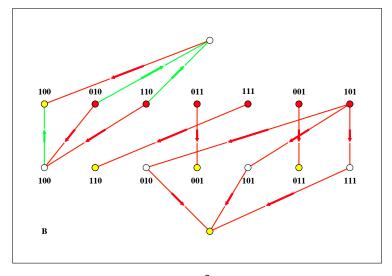
Ordering on 1-dimensional subspaces ??

Internal/external activity induces partition of lattice in prime-free minors; that gives the Tutte polynomial.

classical: every part contains a basis

q: several bases per part, what is the right partition?

So the q-Tutte polynomial is a sum over parts of the partition: exponents of x and y depend on rank/nullity of the parts.



 $T(x,y) = x^2 + xy + 3x$

What's next?

Work in progress:

- ► q-analogue of Tutte polynomial
- ► Link with rank weight enumerator
- ▶ Do all q-matroids come from rank metric codes? How?

Long term:

- ► More cryptomorphic descriptions (circuits, flats, closure, ...)
- ▶ Rank metric codes that are not \mathbb{F}_{q^m} -linear
- Puncturing and shortening of rank metric codes vs. restriction and contraction of q-matroids?
- ► Link with other *q*-analogues?